Peer-to-peer Affine Commitment using Bitcoin

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June 17, 2015

Massively Multiplayer Online Linear Logic

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- A general peer-to-peer commitment mechanism using the language of linear logic
- ► Implemented on top of the Bitcoin network
- ► With applications for proof-carrying authorization

 ${\it Proof-carrying}\ authorization$

 Idea: represent authorization as logical propositions (Appel and Felten 1999)

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- ▶ ... in a logic with a notion of affirmation
- $\langle K \rangle A$ means "the principal K says A"

Proof-carrying authorization

► Alice wants to give access to a file, so affirms:

- ► ⟨Alice⟩ may-read(Charlie, file)





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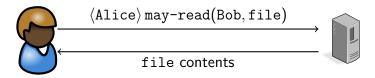
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Proof-carrying authorization - higher order use

• Much more flexible policies are possible:

 $\langle \text{Alice} \rangle \forall K. \langle \text{Registrar} \rangle \text{ in-Alice's-class}(K) \\ \supset \text{may-read}(K, \text{file})$

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Then can derive:

$$orall K. \langle \operatorname{Registrar} \rangle \operatorname{in-Alice's-class}(K) \ \supset \langle \operatorname{Alice} \rangle \operatorname{may-read}(K, \operatorname{file})$$

 $Implementing\ proof-carrying\ authorization$

- Straightforward to make work even in a decentralized/peer-to-peer system
- Proofs are self-contained
- Digital signatures used for affirmation

 $Consumable\ credentials$

What if we want one time use authorization?



Linear logic

- ► Garg et al. 2006; *linear* proof-carrying authorization
- Linear logic treats hypotheses as scarce resources that must be used once

For logicians

Linear logic allows exchange, but not weakening or contraction

Linear logic

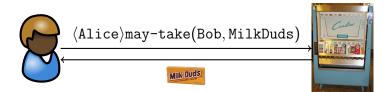
- ► Garg et al. 2006; *linear* proof-carrying authorization
- Linear logic treats hypotheses as scarce resources that must be used once
- Good for modeling state change:

bread \otimes ham \multimap ham sandwich $\forall i. \text{ counter}(i) \multimap \text{ counter}(i+1)$

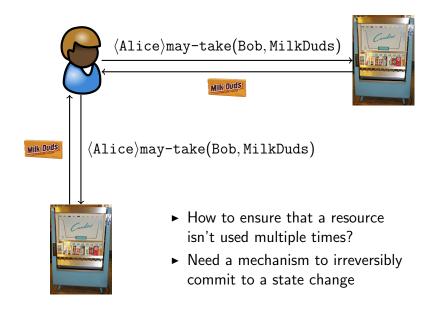
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 $Linear\ authorization$

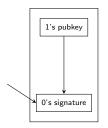


 $Linear\ authorization$

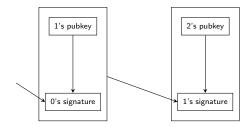


 On a completely different note: consider designing a decentralized digital currency

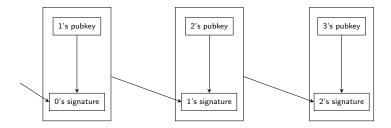
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- A coin is a chain of digital certificates
- ► A coin is spent by signing it over to somebody else

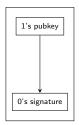


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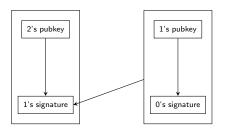


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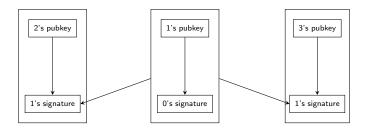




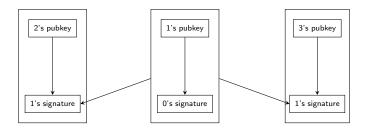
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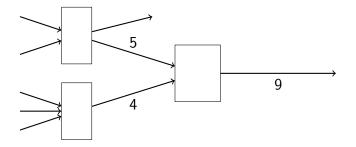


- But how do we prevent an owner from spending a coin multiple times?
- ► Need a mechanism to irreversibly commit to a state change

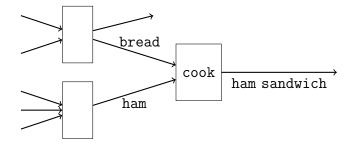
Bitcoin implementation

- Bitcoin (Nakamoto 2008) does this with a global ledger of all transactions - the "blockchain"
- ► Ledger maintained by distributed process called "mining"

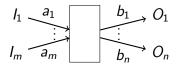
From Bitcoin to Typecoin



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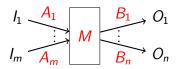


From Bitcoin to Typecoin - transactions



$$\bullet \ a_1 + \cdots + a_m = b_1 + \cdots + b_n$$

From Bitcoin to Typecoin - transactions



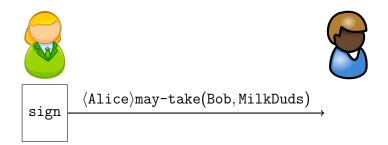
- $\blacktriangleright \vdash M: (A_1 \otimes \cdots \otimes A_m) \multimap (B_1 \otimes \cdots \otimes B_n)$
- Carry linear logic¹ propositions instead of numbers

¹actually affine logic

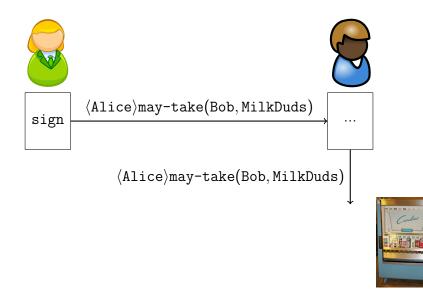


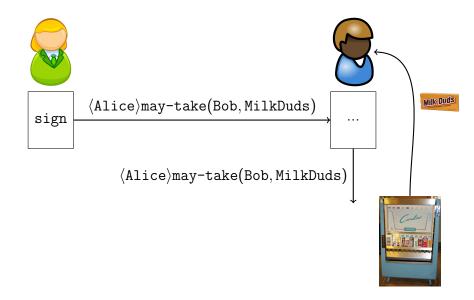


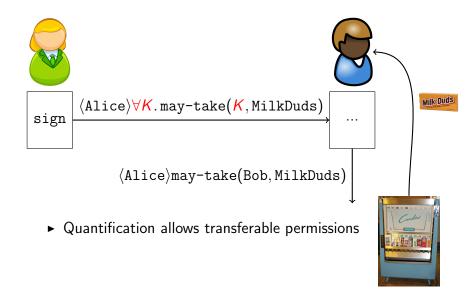


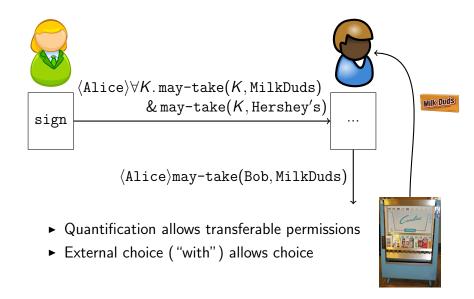












Declarations

Where do may-take, MilkDuds, etc. come from?

Declarations

- ▶ Where do may-take, MilkDuds, etc. come from?
- Transactions can declare types and propositions

 $\texttt{may-take} \quad : \quad \texttt{principal} \rightarrow \texttt{candy} \rightarrow \texttt{prop}$

Building a new currency

Can turn Typecoin back into a currency (S-coins)

 $\begin{array}{rcl} \operatorname{coin} & : & \operatorname{nat} \to \operatorname{prop} \\ \operatorname{merge} & : & \forall \, N, \, M : \operatorname{nat.} \\ & & \operatorname{coin} N \otimes \operatorname{coin} M \multimap \operatorname{coin} N + M \\ \operatorname{split} & : & \forall \, N, \, M, \, P : \operatorname{nat.} \\ & & \operatorname{coin} N + M \multimap \operatorname{coin} N \otimes \operatorname{coin} M \end{array}$

Central banking

Need some way to mint a new S-coin

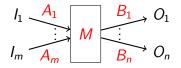
print : nat \rightarrow prop issue : $\forall N:$ nat. $\langle \text{Janet} \rangle$ (print N) \multimap coin N

- ► We *could* build Typecoin in a standalone way
- Use adapted versions of the Bitcoin mining algorithms and protocol
- ► Could typecheck transactions before they enter the chain

- ▶ How to incentivize people to mine on a Typecoin chain?
- Bitcoin already has a lot of mining power
- Typechecking transactions in the chain not an obvious win: proofs might be big or not public

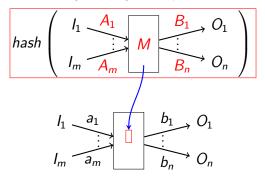
Overlaying on Bitcoin

New plan: actually overlay on top of Bitcoin



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- Embed a hash in the metadata of the Bitcoin transaction
- Send the Typecoin transactions to interested parties

 $Metadata\ in\ Bitcoin$

- Bitcoin historically lacked a nice place to put metadata on principle
- (Nodes would not forward transactions that used the straightforward methods)
- Paper describes a somewhat hacky workaround

Metadata in Bitcoin

- Bitcoin historically lacked a nice place to put metadata on principle
- (Nodes would not forward transactions that used the straightforward methods)
- Paper describes a somewhat hacky workaround
- But the Bitcoin developers have since caved

• Receipts that attest to outputs: receipt($A \rightarrow addr$)

$$\langle \text{Alice} \rangle (\text{receipt}(\text{coin}(5) \rightarrow \text{Alice}) \rightarrow \forall K. \text{may-take}(K, \text{MilkDuds}))$$

Expiration/revocation

 Conditional modality permits revocation and expiration: if(before(July 10)), may-write(Alice, POPL-paper))

$\ Implementation$

- Implemented in Standard ML
- ▶ With a new Bitcoin client, in SML

$Related \ Work$

- ► Bowers et al. 2007; consumable credentials
- ► Rosenfeld 2013; colored coins
- ► Wood 2014; Ethereum

Conclusion

- Typecoin is a flexible peer-to-peer logical commitment mechanism
- Based on generalizing Bitcoin to carry logical propositions
- Actually implemented on top of Bitcoin
- Details on the logic are in the paper

Thank you!

Why not linear?

- Typecoin sort of *fundamentally* affine can always throw away an output
- ► Allowing rule declarations in signatures makes it trivial
- ▶ trash : ⊤ → 1

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- ► Allowing rule declarations in signatures makes it trivial
- trash : $\top \multimap 1$
- Prohibit \top ? trash : A \multimap 1
- ▶ Prohibit proving 1? dummy : prop. trash : A !dummy
- ► Prohibit consuming A? trash : ⟨K⟩dummy --∞ !dummy, sign ⟨K⟩(A --∞ dummy)

Metadata: "m-of-n" outputs

- ► An "*m*-of-*n*" output lists *n* public keys
- ► To spend it, provide signatures using *m*
- ► 2-of-3 outputs useful for two-party escrow
- ► We use 1-of-2 outputs to embed metadata
- One public key is the real destination
- ► The other is actually the hash of our transaction