

# *Closures for Rust*

Michael Sullivan

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# *Outline*

*Introduction*

*Rust*

*Closures*



## *Disclaimer*



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- What I discuss and how I present issues reflect my personal biases in language design.



## *Goals*

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*What do we want in a programming language?*

- Fast: generates efficient machine code
- Safe: type system provides guarantees that prevent certain bugs
- Concurrent: easy to build concurrent programs and to take advantage of parallelism
- “Systemsy”: fine grained control, predictable performance characteristics

*Goals*  
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- Firefox is in C++, which is Fast and Systemsy
- ML is (sometimes) fast and (very) safe
- Erlang is safe and concurrent
- Haskell is (sometimes) fast, (very) safe, and concurrent
- Java and C# are fast and safe





# *Rust*

a systems language  
pursuing the trifecta  
safe, concurrent, fast  
-lkuper



# *Rust*

Design  
Status



*Design*  
*Goals (straight from the docs)*

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*Design*  
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- Compile-time error detection and prevention
- Run-time fault tolerance and containment
- System building, analysis and maintenance affordances
- Clarity and precision of expression
- Implementation simplicity
- Run-time efficiency
- High concurrency

*Design*  
*Type system features*

- Algebraic data type and pattern matching (no null pointers!)
- Polymorphism: functions and types can have generic type parameters
- Type inference on local variables
- Lightweight object system
- Data structures are immutable by default

*Design*  
*Other features*

- Lightweight tasks with no shared state
- Control over memory allocation
- Move semantics, unique pointers
- Function arguments can be passed by alias
- Typestate system tracks predicates that hold at points in the program

*Design*  
*... What?*

“It’s like C++ grew up, went to grad school, started dating ML, and is sharing an office with Erlang.”



# *Status rustc*

- Self-hosting rust compiler



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### *rustc*

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## *Status*

### *rustc*

- Self-hosting rust compiler
- Uses LLVM as a backend
- Handles polymorphism through type passing (blech)
- Memory management through automatic reference counting (eww)



*Status*  
*The catch*

- Not ready for prime time



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- Lots of bugs and exposed sharp edges



*Status*  
*The catch*

- Not ready for prime time
- Lots of bugs and exposed sharp edges
- Language still changing rapidly

## *Closures*

What closures are  
Closures in rust

*What closures are*  
*Definition*

- In civilized languages, functions are first-class values and are allowed to reference variables in enclosing scopes
- That is, they close over their environments

*What closures are*  
*Example*

```
function add(x) {  
  return function(y) { return x + y; };  
}  
var foo = add(42)(1337); // 1379
```

- Produces a function that adds x to its argument
- Note that the inner function outlives the enclosing function. x can't just be stored on the stack.

*What closures are*  
*Another Example*

```
function scale(x, v) {  
  return map(function(y) { return x * y; },  
             v);  
}  
var v = scale(2, [1, 2, 3]); // [2, 4, 6]
```

- Multiplies every element in an array by some amount
- Note that here the lifetime of the inner function is shorter than the lifetime of the enclosing one. `x` could just be stored on the stack.



*What closures are*  
*Traditional implementation*

- Represent functions as a code pointer, environment pointer pair

*What closures are*  
*Traditional implementation*

- Represent functions as a code pointer, environment pointer pair
- Heap allocate stack frames (or at least the parts that are closed over)



*Closures in rust*  
*Design constraints*

- Want to be explicit about when we are allocating memory
- Don't want to have to heap allocate closures when it isn't necessary



## *Closures in rust*

### *Solutions*

- Have two function types: **block** and **fn**

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- Have two function types: **block** and **fn**
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- Values of **fn** type can be automatically coerced to **block** type when passed as function arguments; this allows more code reuse

*Closures in rust*  
*Solutions*

- Have two function types: **block** and **fn**
- Values of a **block** type may not be copied, but can be passed by alias; this prevents them from escaping
- Values of **fn** type can be automatically coerced to **block** type when passed as function arguments; this allows more code reuse
- Explicitly state what sort of function you writing

## *Closures in rust*

### *Lambda example*

```
fn add(x: int) -> fn(int) -> int {  
  ret lambda(y: int) -> int { ret x + y; };  
}
```

- lambda produces a **fn** that closes over its environment by copying upvars into a heap allocated environment
- Since the variables are copied, changes made to the variables in the enclosing scope will not be reflected in the nested function



## *Closures in rust*

### *Block example*

```
fn scale(x: int, v: &[int]) -> [int] {  
  map(block(y: &int) -> int { x * y }, v)  
}
```

- block produces a **block** that closes over its environment by storing pointers to the stack locations of the variables in a stack allocated environment

*Closures in rust*  
*Inference example*

```
fn scale(x: int, v: &[int]) -> [int] {  
    map(|&y| x * y), v)  
}
```

- Provides an abbreviation for block; the argument and return types are type inferred
- Only allowed to appear as a function argument, making type inference easy

*Closures in rust*  
*Coercion example*

```
fn add1(x: &int) -> int { x + 1 }  
fn increment(v: &[int]) -> [int] {  
    map(add1, v)  
}
```

- add1 is coerced to a **block** when passed to map

## *Conclusion*

- Rust is a new systems language out of Mozilla Research that is designed to be fast, concurrent, and safe
- Closures are a tricky design space in languages that want to be explicit about performance
- Rust approaches the issues by separating functions into multiple varieties.